

## wp Calculation for Sequential Composition

## Correctness of Programs: Sequential Composition

Is  $\{ \text{True} \} \text{tmp} := x; x := y; y := \text{tmp} \{ x > y \}$  correct?

## Rules of Weakest Precondition: Summary

$$wp(x := e, R) = R[x := e]$$

$$wp(\text{if } B \text{ then } S_1 \text{ else } S_2 \text{ end}, R) = \left( \begin{array}{c} B \Rightarrow wp(S_1, R) \\ \wedge \\ \neg B \Rightarrow wp(S_2, R) \end{array} \right)$$

$$wp(S_1 ; S_2, R) = wp(S_1, wp(S_2, R))$$

# Proof Rules using Weakest Precondition

$$\{Q\} S \{R\} \equiv Q \Rightarrow wp(S, R)$$

$$\{Q\} x := e \{R\} \iff Q \Rightarrow \underbrace{R[x := e]}_{wp(x := e, R)}$$

$$\begin{aligned} & \{Q\} \text{ if } B \text{ then } S_1 \text{ else } S_2 \text{ end } \{R\} \\ & \iff \left( \begin{array}{c} \{Q \wedge B\} S_1 \{R\} \\ \wedge \\ \{Q \wedge \neg B\} S_2 \{R\} \end{array} \right) \iff \left( \begin{array}{c} (Q \wedge B) \Rightarrow wp(S_1, R) \\ \wedge \\ (Q \wedge \neg B) \Rightarrow wp(S_2, R) \end{array} \right) \end{aligned}$$

$$\{Q\} S_1 ; S_2 \{R\} \iff Q \Rightarrow \underbrace{wp(S_1, wp(S_2, R))}_{wp(S_1 ; S_2, R)}$$

# Correctness of Loops

```
{ Q }
```

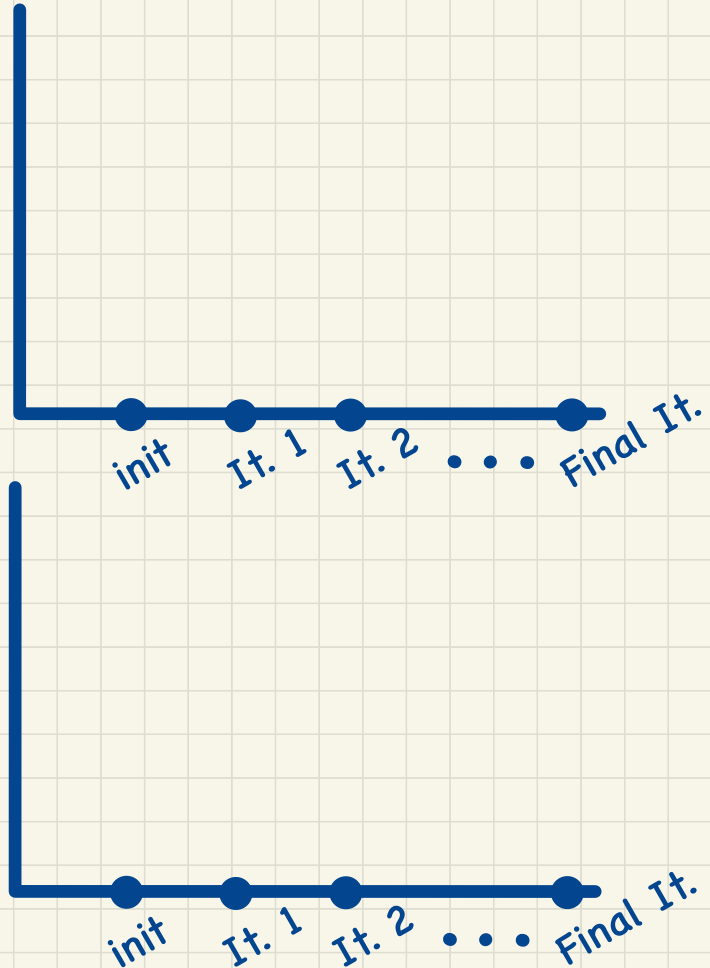
```
  Sinit
```

```
  while ( B ) {
```

```
    Sbody
```

```
  }
```

```
{ R }
```



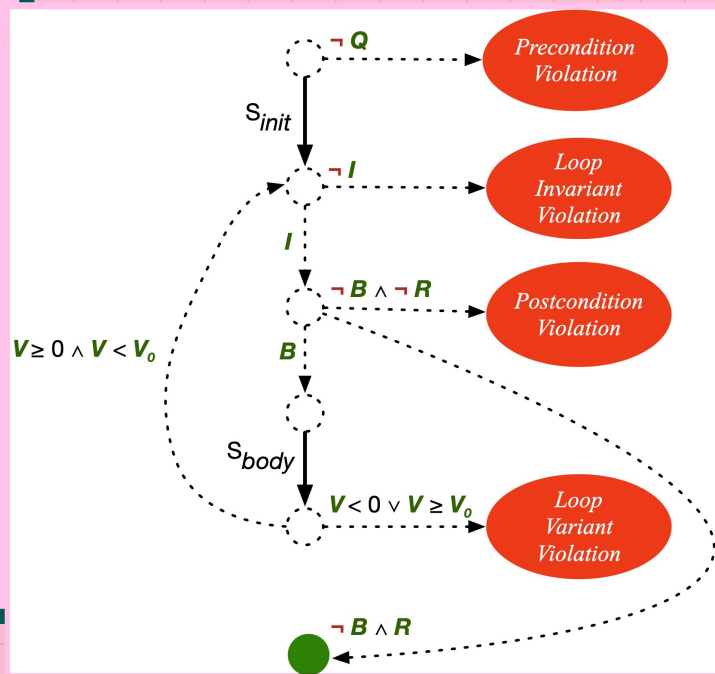
# Contracts of Loops

## Syntax

```
CONSTANT ... (* input list *)
I(var_list) == ...
V(var_list) == ...
--algorithm MYALGORITHM {
  variables ..., variant_pre = 0, variant_post = 0
  {
    assert Q; (* Precondition *)
    S_init
    assert I(...); (* Is LI established? *)
    while( B ) {
      variant_pre := V(...);
      S_body
      variant_post := V(...);

      assert variant_post >= 0;
      assert variant_post < variant_pre;
      assert I(...); (* Is LI preserved? *)
    }
    assert R; (* Postcondition *)
  }
}
```

## Runtime Checks



## Contracts of Loops: Example

Assume: Q and R are true

```

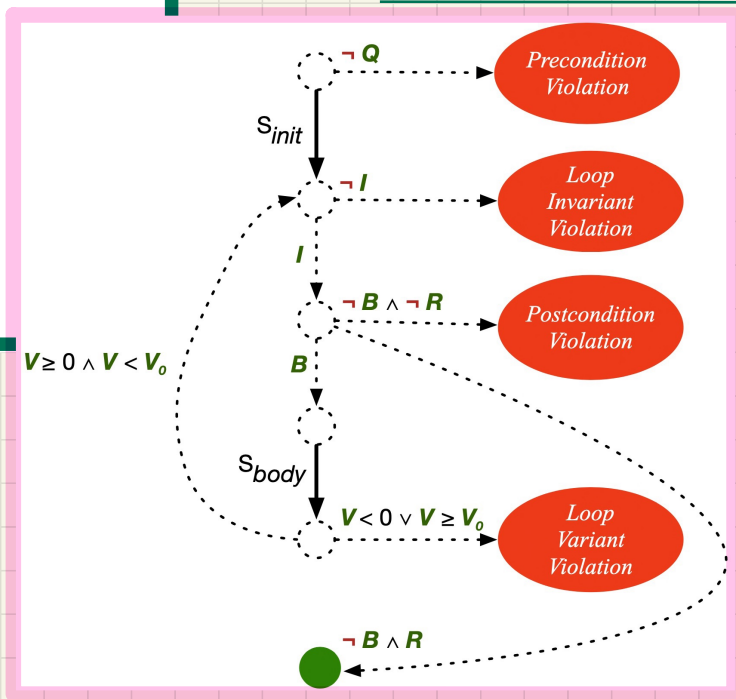
1  I(i) == (1 <= i) /\ (i <= 6)
2  V(i) == 6 - i
3  --algorithm loop_invariant_test
4    variables i = 1, variant_pre = 0, variant_post = 0;
5    {
6      assert I(i);
7      while (i <= 5) {
8        variant_pre := V(i);
9        i := i + 1;
10       variant_post := V(i);
11       assert variant_post >= 0;
12       assert variant_post < variant_pre;
13       assert I(i);
14     } ;
15   }

```

## Specification

## Specification

# Runtime Checks

[illegible]

# Contracts of Loops: Violations

Assume: Q and R are **true**

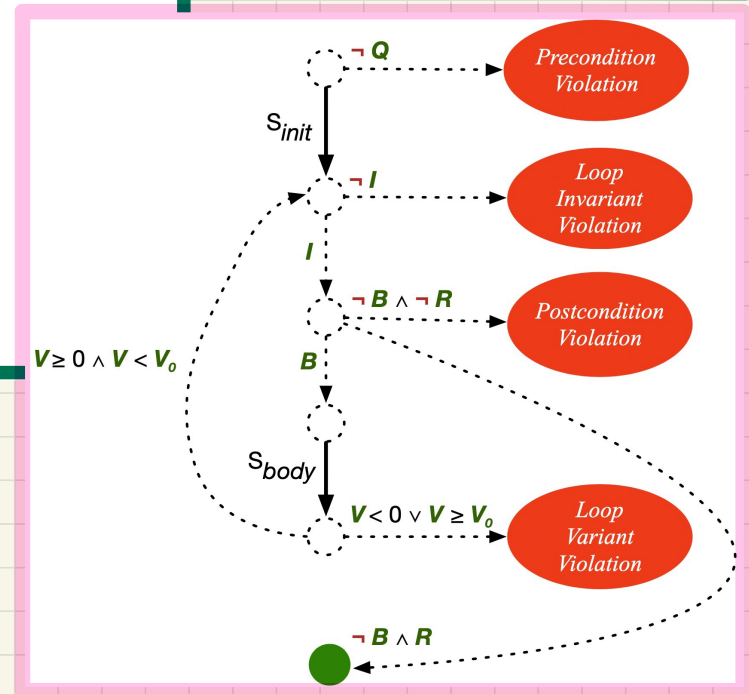
```
1  I(i) == (1 <= i) /\ (i <= 6)
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3  --algorithm loop_invariant_test
4    variables i = 1, variant_pre = 0, variant_post = 0;
5    {
6      assert I(i);
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8        variant_pre := V(i);
9        i := i + 1;
10       variant_post := V(i);
11       assert variant_post >= 0;
12       assert variant_post < variant_pre;
13       assert I(i);
14     } ;
15 }
```

## Specification

## Runtime Checks

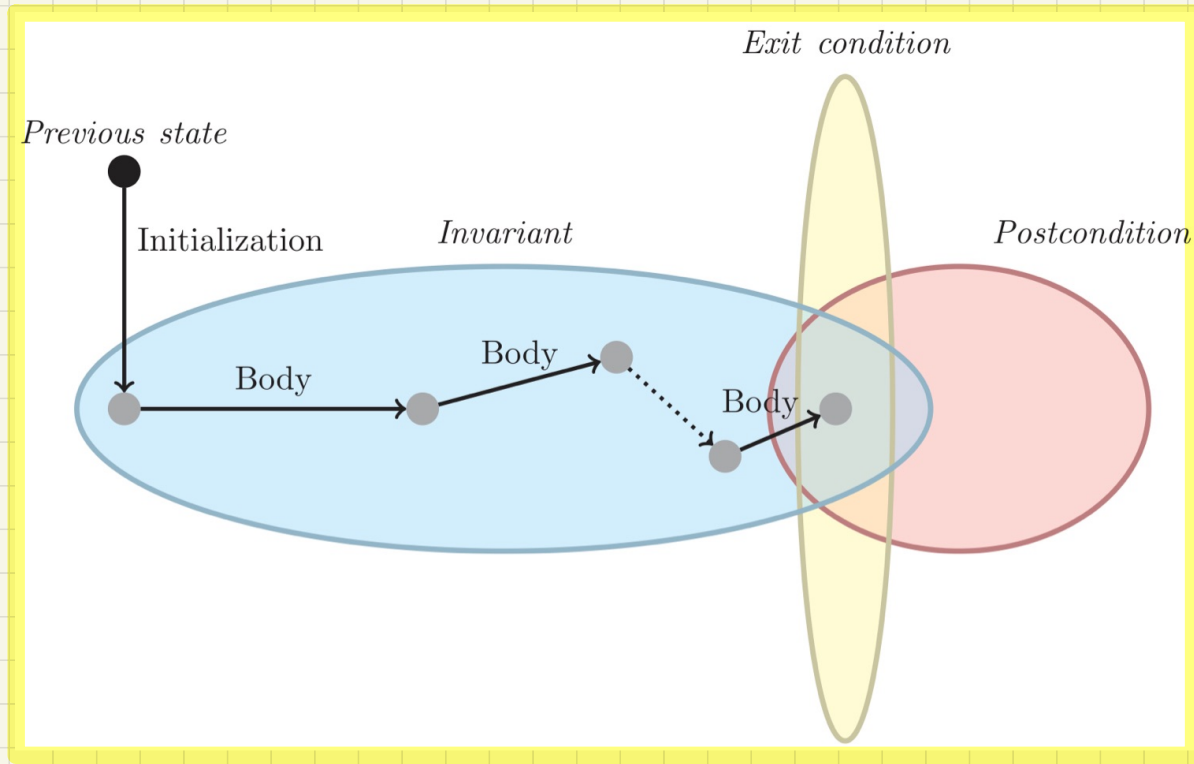
invariant:  $1 \leq i \leq 5$

variant:  $5 - i$





# Contracts of Loops: Visualization



# Correct Loops: Proof Obligations

```
{Q}
Sinit
assert I(...);
while( B ) {
  variant_pre := V(...);
  Sbody
  variant_post := V(...);
  assert variant_post >= 0;
  assert variant_post < variant_pre;
  assert I(...);
}
{R}
```

- A loop is *partially correct* if:
  - Given precondition  $Q$ , the initialization step  $S_{init}$  establishes  $LI$ .
  - At the end of  $S_{body}$ , if not yet to exit,  $LI$  is maintained.
  - If ready to exit and  $LI$  maintained, postcondition  $R$  is established.
- A loop *terminates* if:
  - Given  $LI$ , and not yet to exit,  $S_{body}$  maintains  $LV$   $V$  as non-negative.
  - Given  $LI$ , and not yet to exit,  $S_{body}$  decrements  $LV$   $V$ .

# Correct Loops: Proof Obligations

## Example

### Specification

```
1  I(i) == (1 <= i) /\ (i <= 6)
2  V(i) == 6 - i
3  --algorithm loop_invariant_test
4    variables i = 1, variant_pre = 0, variant_post = 0;
5    {
6      assert I(i);
7      while (i <= 5) {
8        variant_pre := V(i);
9        i := i + 1;
10       variant_post := V(i);
11       assert variant_post >= 0;
12       assert variant_post < variant_pre;
13       assert I(i);
14     } ;
15 }
```

- A loop is **partially correct** if:
  - Given precondition **Q**, the initialization step  $S_{init}$  establishes **LI**  $I$ .
  - At the end of  $S_{body}$ , if not yet to exit, **LI**  $I$  is maintained.
  - If ready to exit and **LI**  $I$  maintained, postcondition **R** is established.
- A loop **terminates** if:
  - Given **LI**  $I$ , and not yet to exit,  $S_{body}$  maintains **LV**  $V$  as non-negative.
  - Given **LI**  $I$ , and not yet to exit,  $S_{body}$  decrements **LV**  $V$ .